

## Exact Minkowski Sums of Polygons With Holes

Alon Baram <sup>1</sup>   Efi Fogel <sup>1</sup>   Dan Halperin <sup>1</sup>  
 Michael Hemmer <sup>2</sup>   Sebastian Morr <sup>2</sup>

<sup>1</sup>



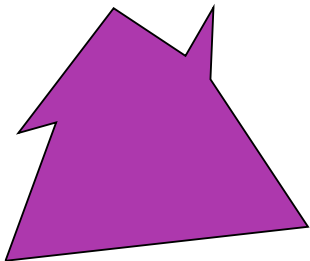
<sup>2</sup>



# What's the Minkowski Sum?

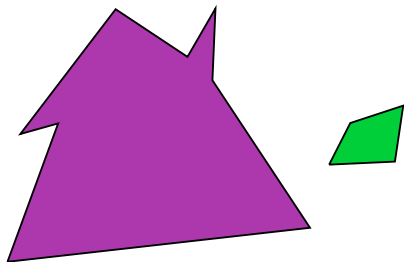
# What's the Minkowski Sum?

$P$



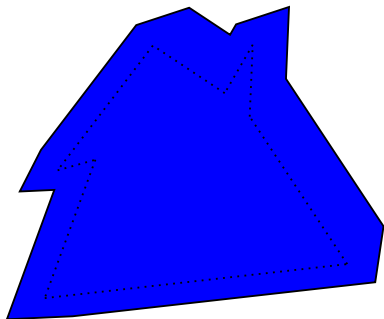
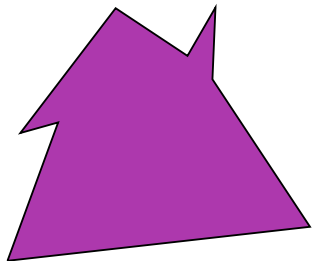
# What's the Minkowski Sum?

$$P \quad \oplus \quad Q$$



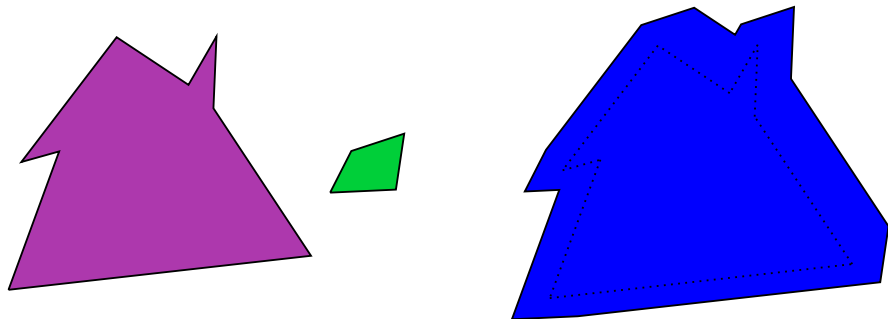
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$$P \oplus Q = \{p + q \mid p \in P, q \in Q\}$$



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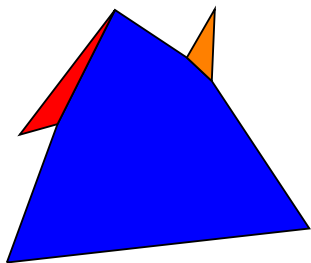


Applications: Motion planning, packing problems, CAD, ...

# Decomposition approach

[Lozano-Pérez (1983)]

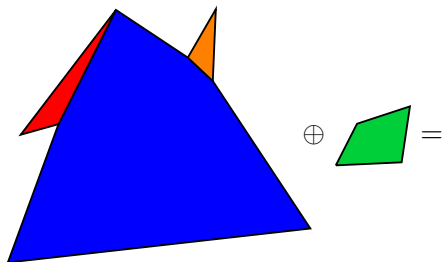
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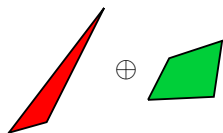
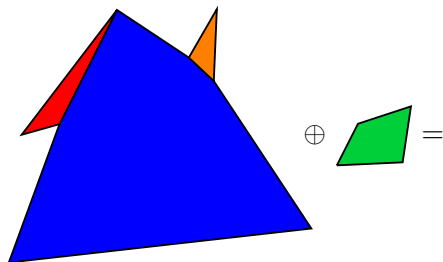


## Decomposition approach



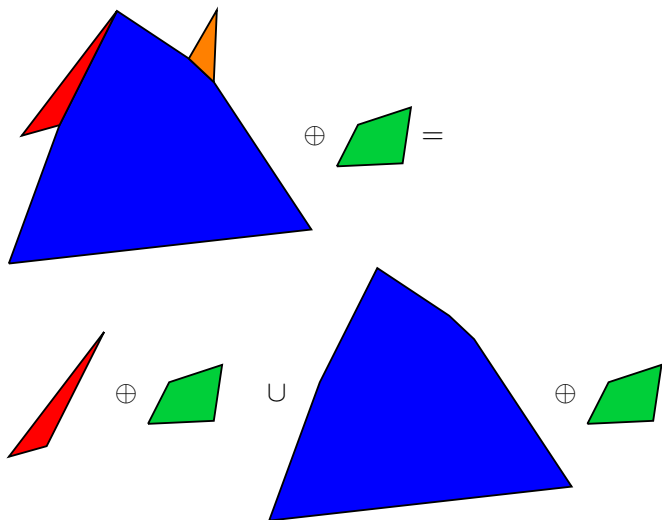
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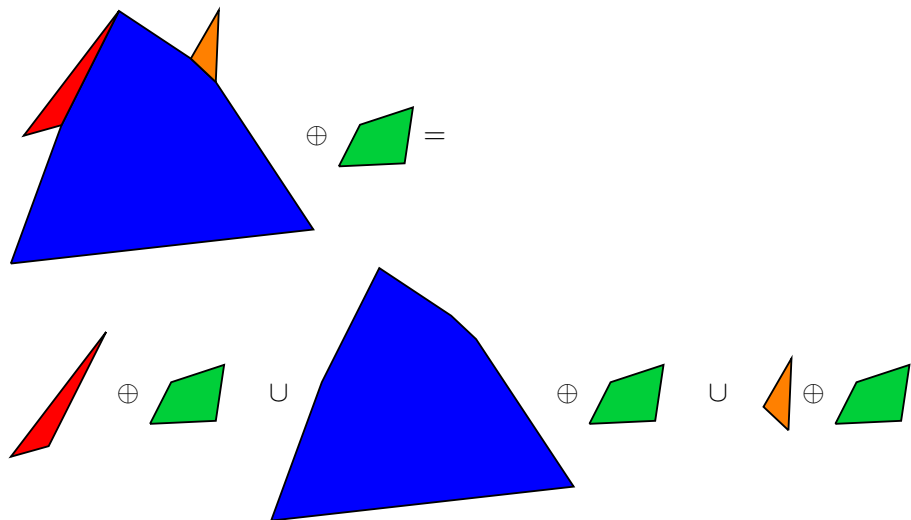
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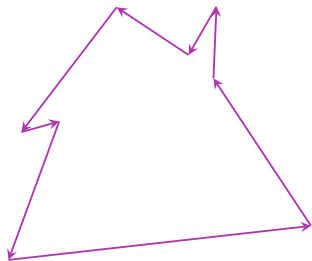
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# Convolution approach

[Guibas et al. (1983)]

## Convolution approach

$$P = (p_0, p_1, \dots, p_{n-1})$$

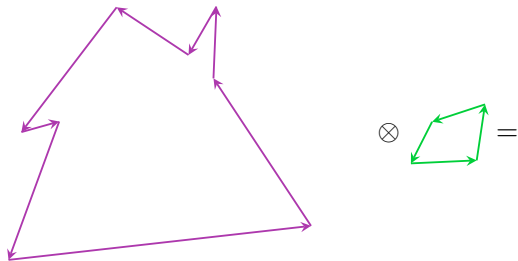


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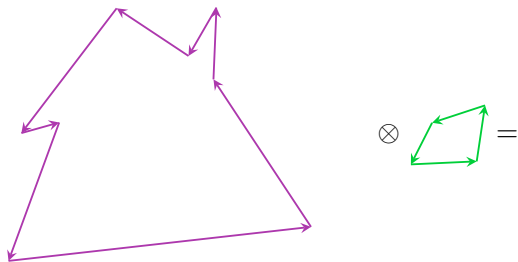
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[Guibas et al. (1983)]

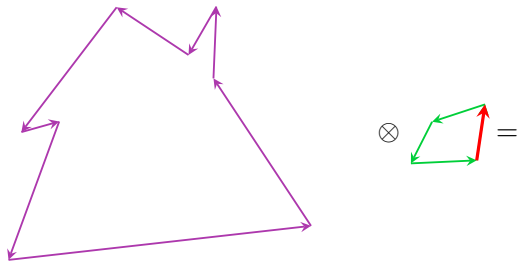


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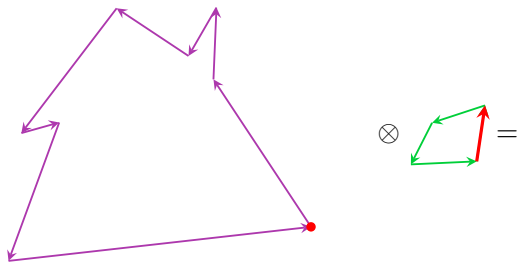
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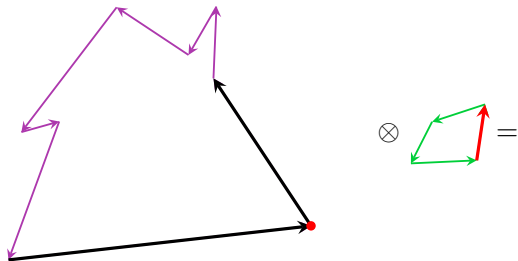
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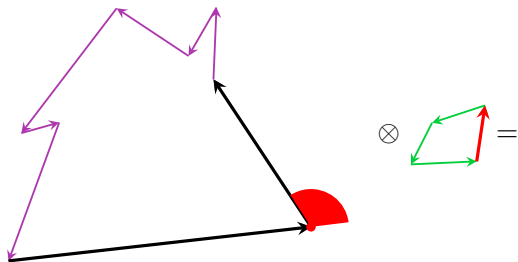
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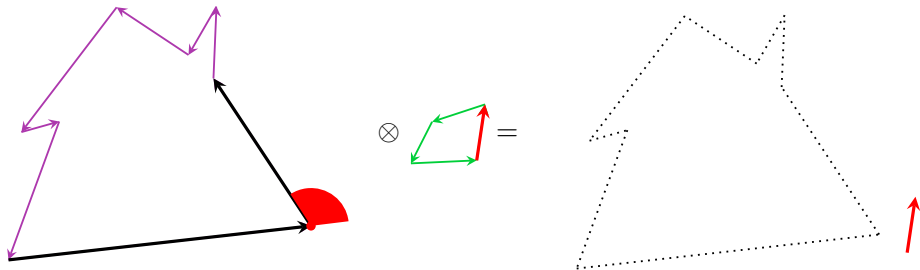
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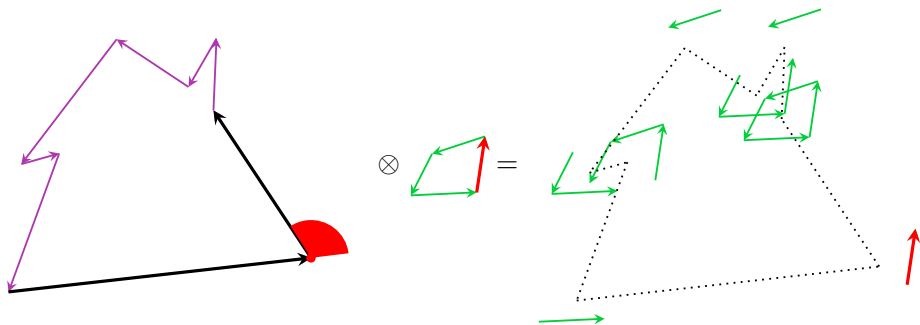
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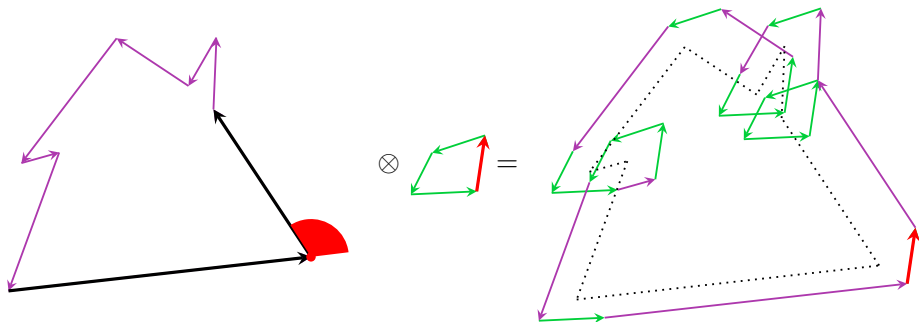
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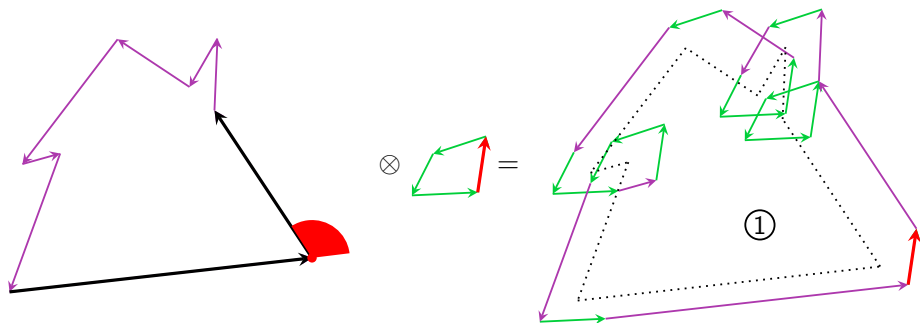
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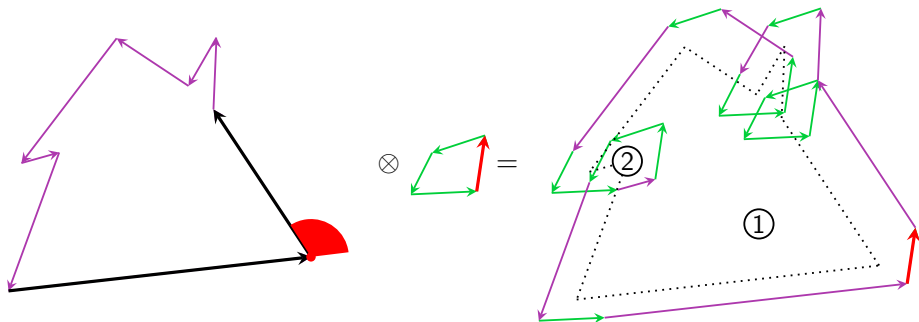


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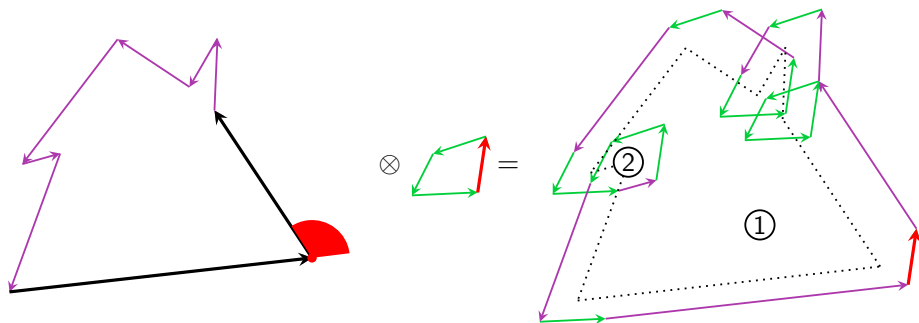
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Bottleneck: arrangement computation, worst case size  $O(n^2 m^2)$

[Guibas et al. (1983)]

# Overview

- 1 Can we speed up the convolution approach?
- 2 Can we fill in holes?
- 3 How does the algorithm compare to other approaches?

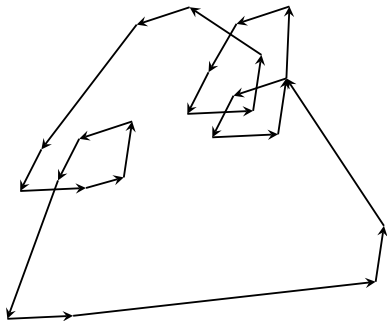
## Observation

Reflex vertices don't contribute to the Minkowski Sum's boundary!

[Kaul et al. (1992)]

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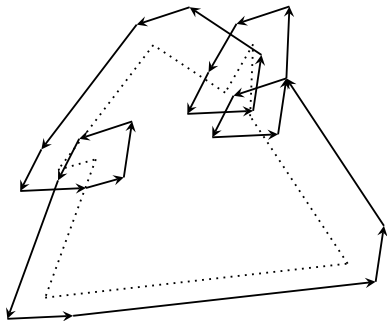
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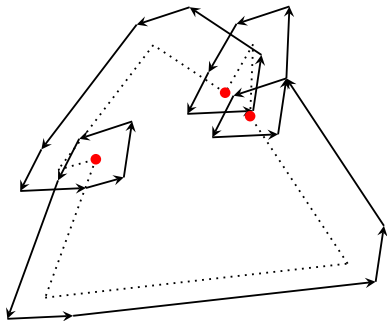
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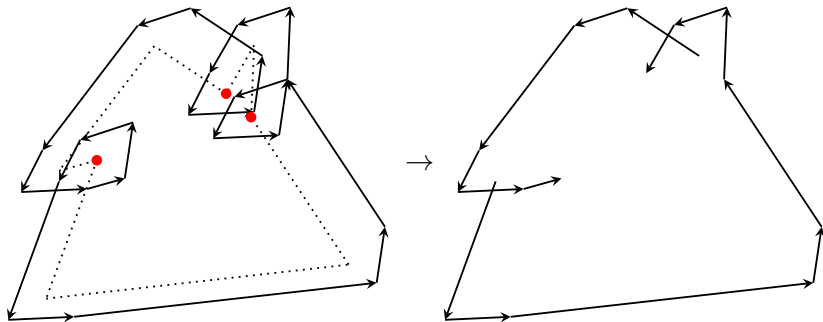
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# Reduced convolution approach

## Step 1

Compute the reduced convolution.

[Behar and Lien, 2011]

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Compute the arrangement of the segments.

The winding number property can not be used anymore.

Instead, two more steps to remove *false holes*.

[Behar and Lien, 2011]

## Quick orientation filter

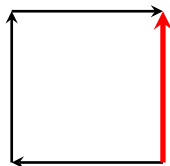
### Step 3

Iterate over all faces, discard loops which are not orientable.

# Quick orientation filter

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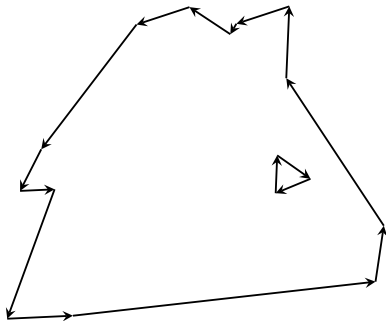
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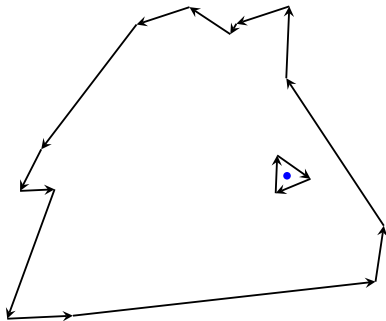
For all remaining faces:





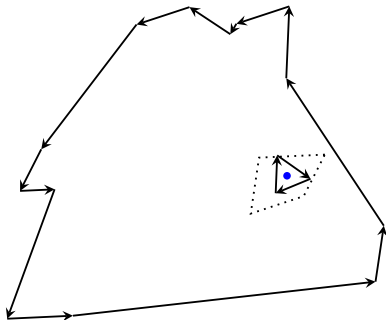
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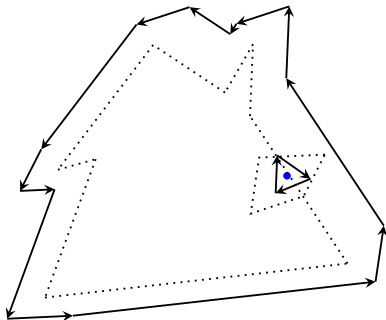
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For all remaining faces: Find a point inside the face, translate  $-Q$  to that point,



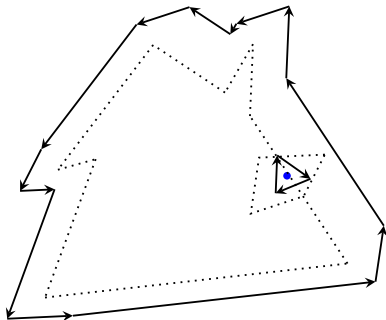
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For all remaining faces: Find a point inside the face, translate  $-Q$  to that point, and intersect it with  $P$ .



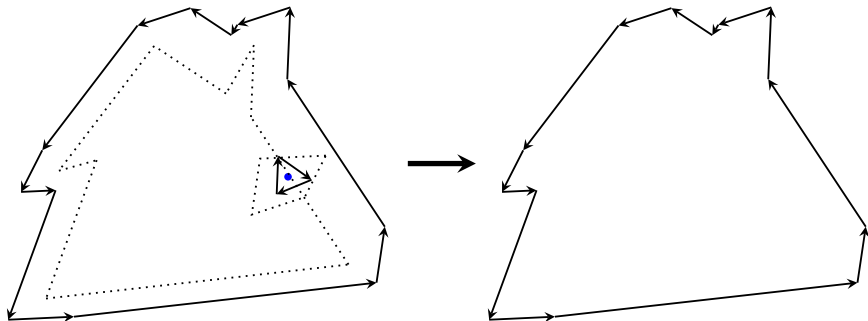
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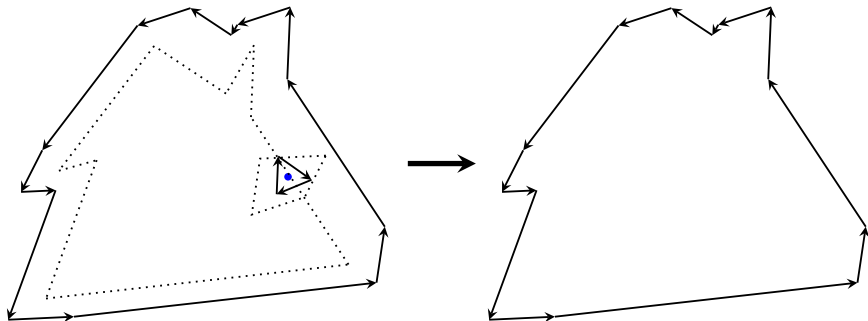
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The remaining faces are exactly the holes of  $P \oplus Q$ !

## What about holes?

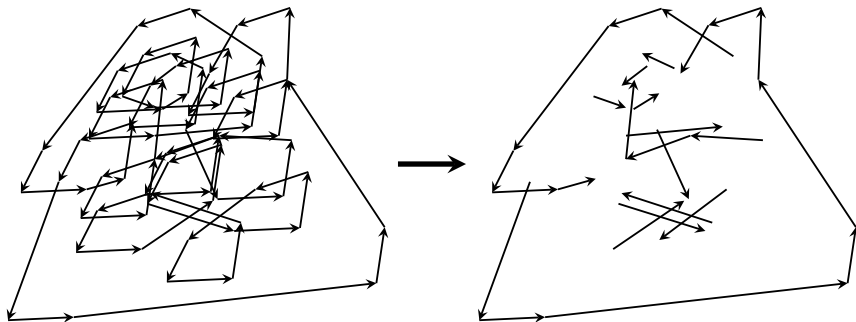
Even the reduced convolution can become quite complex:





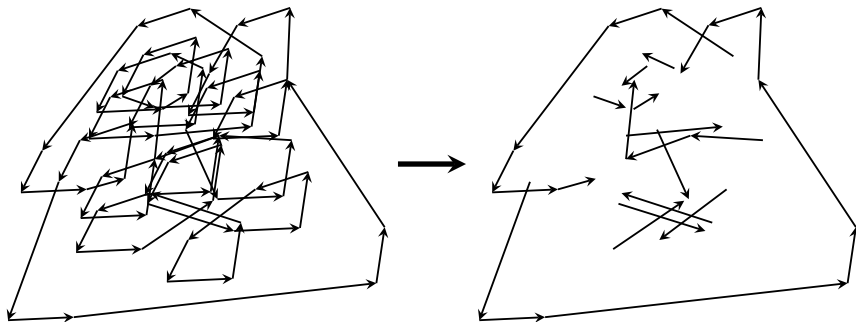
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Idea: Reduce input complexity!

# Overview

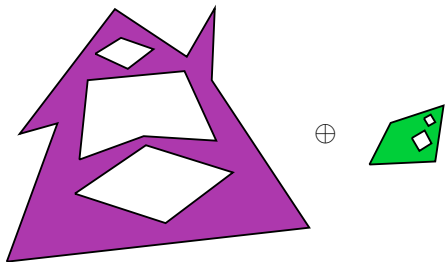
- 1 Can we speed up the convolution approach?
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First, we need some polygons with holes. . .

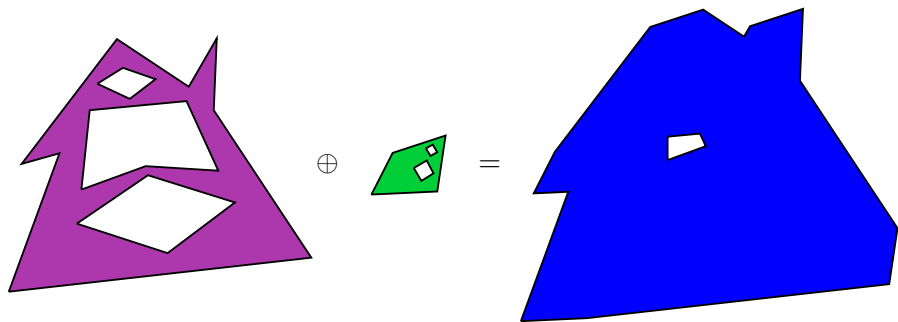
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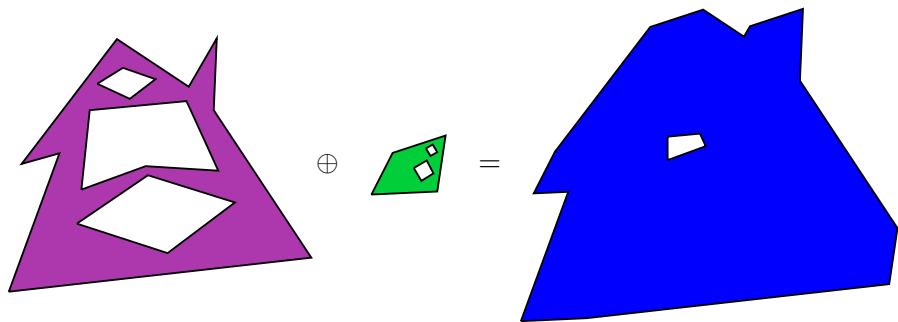
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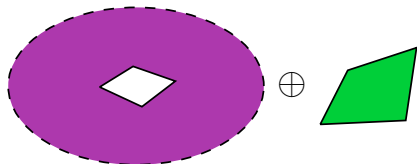


Which holes are relevant?



# Hole filter

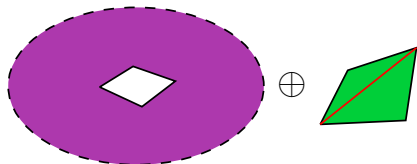
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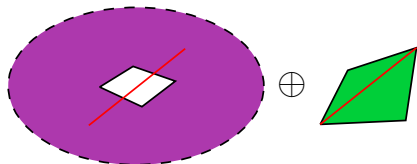
If there is a path  $\gamma$  in  $Q$



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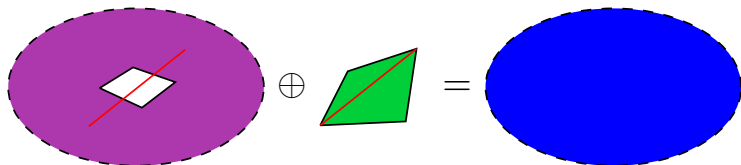
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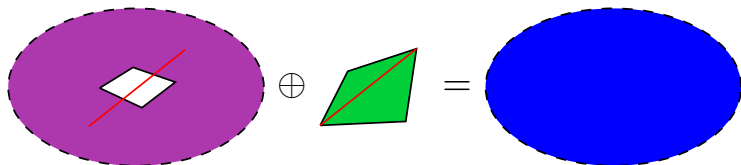
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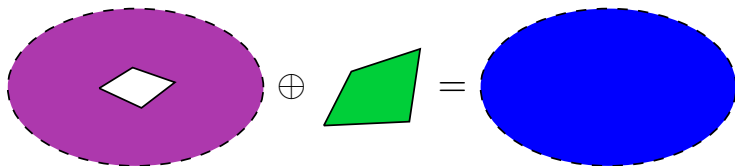
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Because when the hole's boundary is added to  $\gamma$ , it "smears" completely over the hole.

# Corollaries

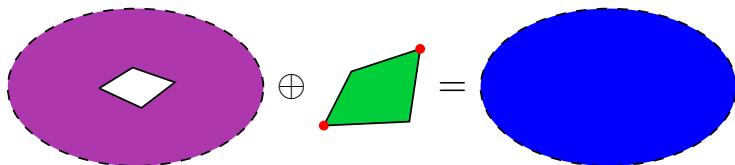
## Corollary



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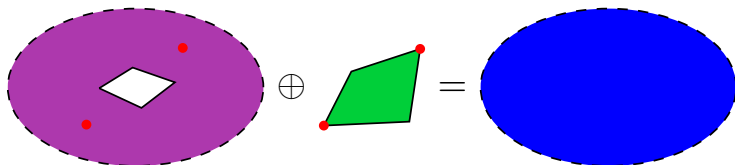
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If there are two points in  $Q$  so that their inverse does not fit under any translation inside a hole of  $P$ ,

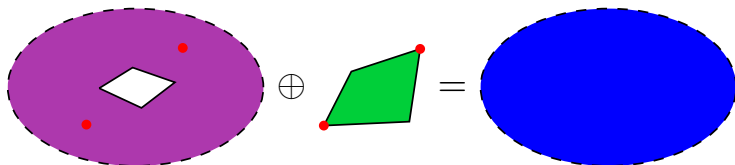




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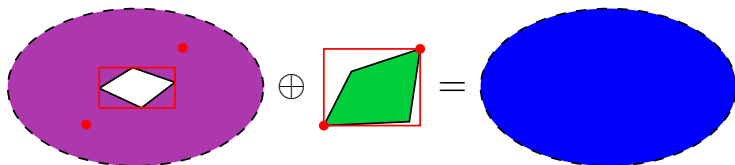
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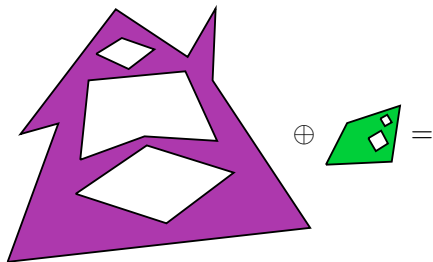
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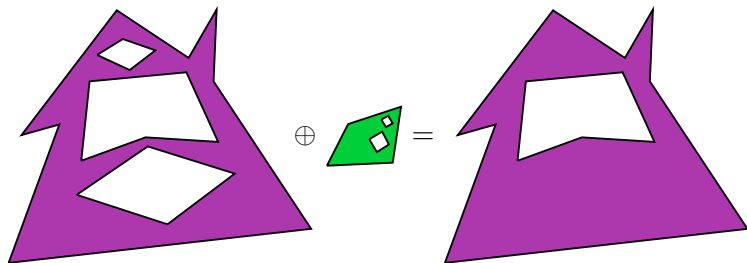
If  $Q$ 's axis-aligned bounding box does not completely fit inside the hole's axis-aligned bounding box, the hole can be filled up.



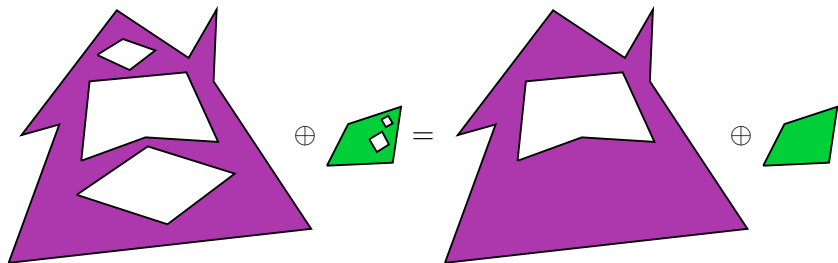
## Effect on input



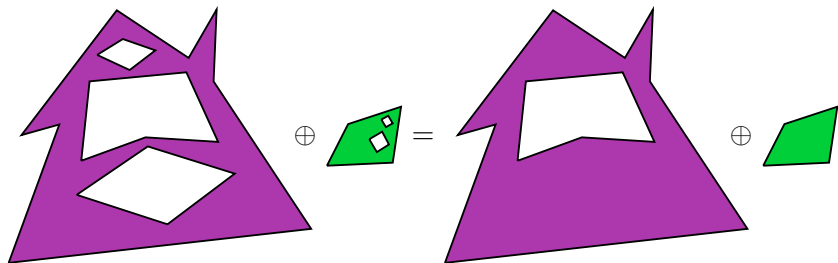
## Effect on input



## Effect on input

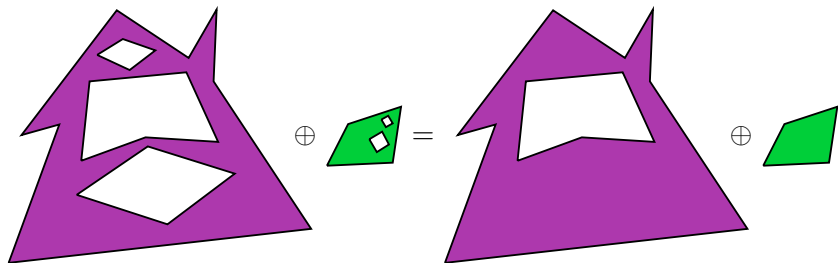


## Effect on input



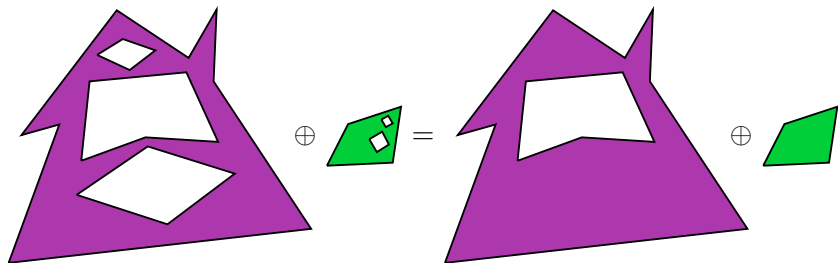
- One is always simple!

## Effect on input



- One is always simple! Sometimes both.

## Effect on input



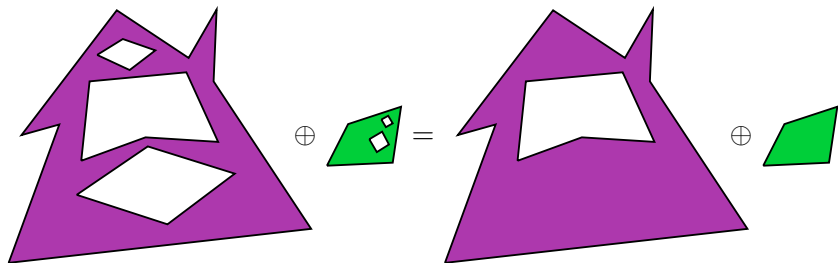
- One is always simple! Sometimes both.

This filter is . . .

- approach independent



## Effect on input



- One is always simple! Sometimes both.

This filter is . . .

- approach independent
- generalizable to higher dimensions

# Overview

- 1 Can we speed up the convolution approach?
- 2 Can we fill in holes?
- 3 How does the algorithm compare to other approaches?

# Implementation

	Simple polygons	General polygons
Inexact		
Exact		

# Implementation

	<b>Simple polygons</b>	<b>General polygons</b>
<b>Inexact</b>		
<b>Exact</b>	Wein (2006)	

# Implementation

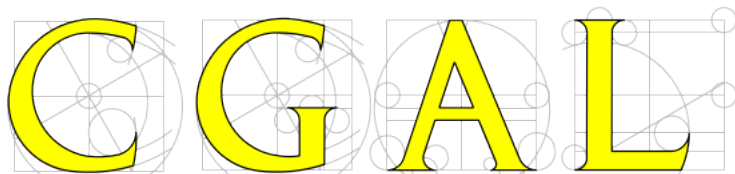
	<b>Simple polygons</b>	<b>General polygons</b>
<b>Inexact</b>		Behar and Lien (2011)
<b>Exact</b>	Wein (2006)	

# Implementation

	<b>Simple polygons</b>	<b>General polygons</b>
<b>Inexact</b>		Behar and Lien (2011)
<b>Exact</b>	Wein (2006)	<b>NEW</b>

# Implementation

	<b>Simple polygons</b>	<b>General polygons</b>
<b>Inexact</b>		Behar and Lien (2011)
<b>Exact</b>	Wein (2006)	<b>NEW</b>



# Implementation

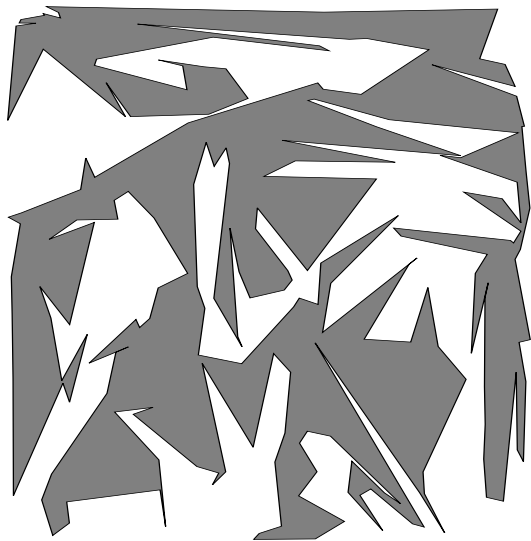
	Simple polygons	General polygons
Inexact		Behar and Lien (2011)
Exact	Wein (2006)	<b>NEW</b>



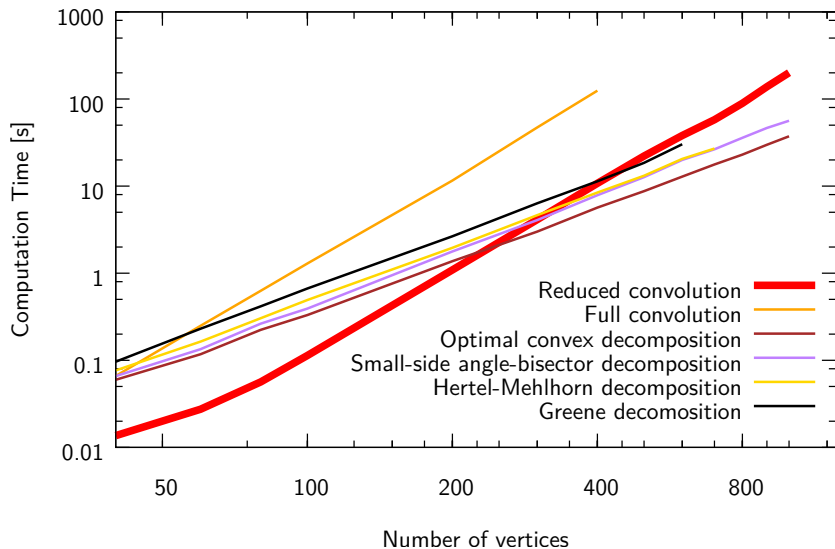
Starting with CGAL 4.7, you can use the method `CGAL::minkowski_sum_2()` on polygons with holes!



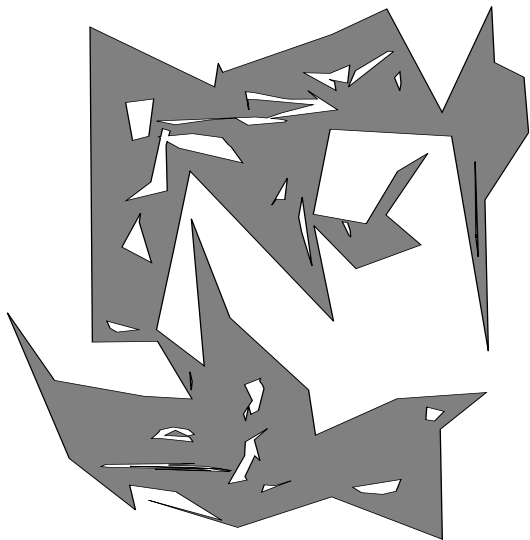
## Benchmark: Simple polygons



# Benchmark: Simple polygons



## Benchmark: Polygons with holes



# Implemented decomposition approaches

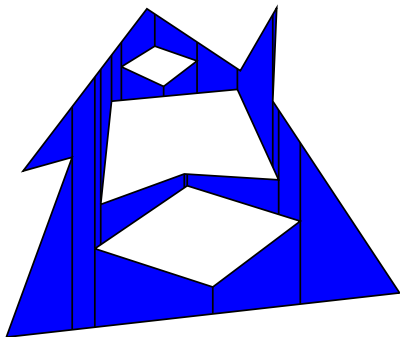
Vertical decomposition

Constrained triangulation

# Implemented decomposition approaches

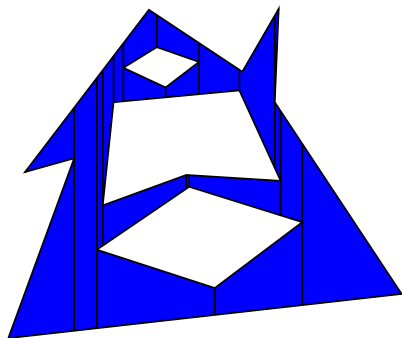
Vertical decomposition

Constrained triangulation

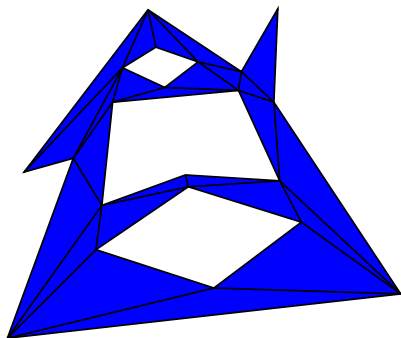


# Implemented decomposition approaches

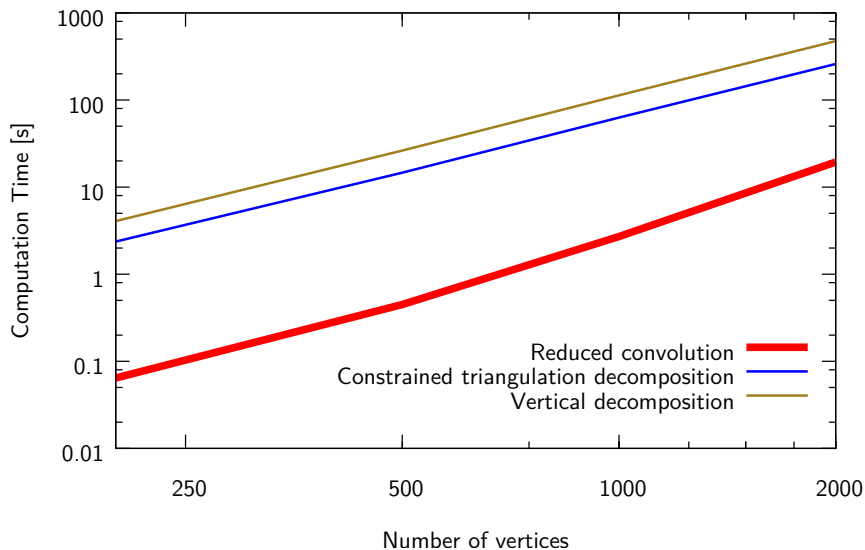
Vertical decomposition



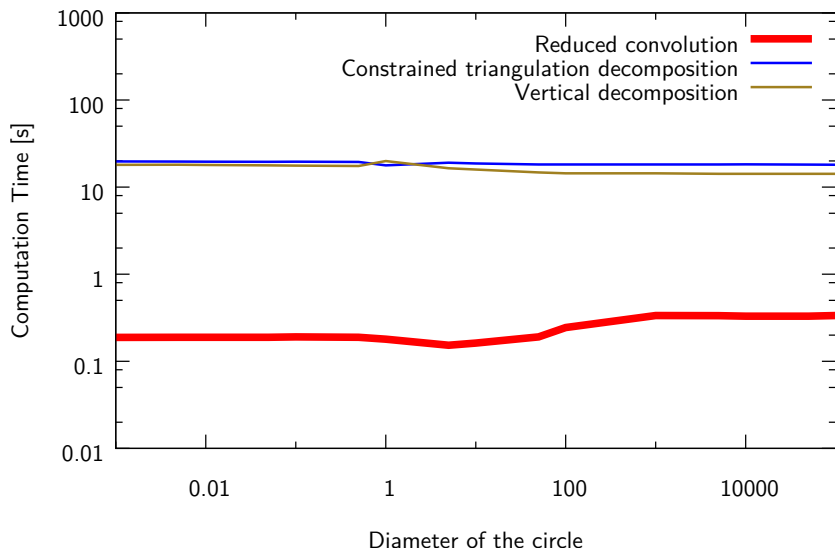
Constrained triangulation



## Benchmark: Polygons with holes

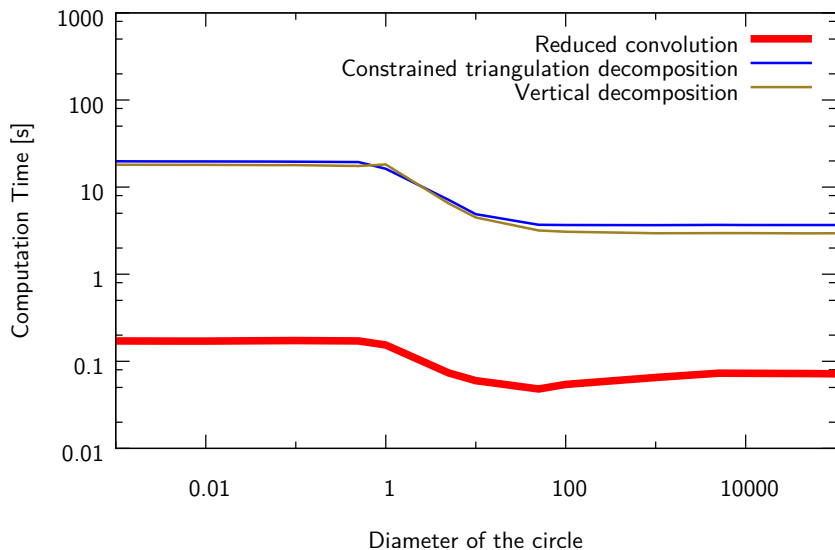


## Benchmark: Growing circle (no hole filter)

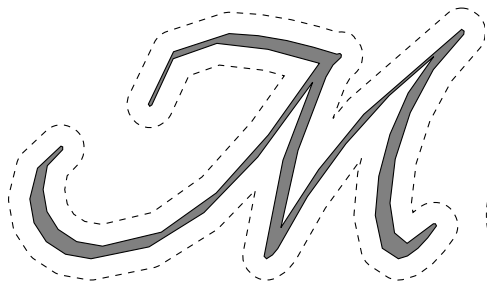




## Benchmark: Growing circle (with hole filter)



## Benchmark: Glyph Offset

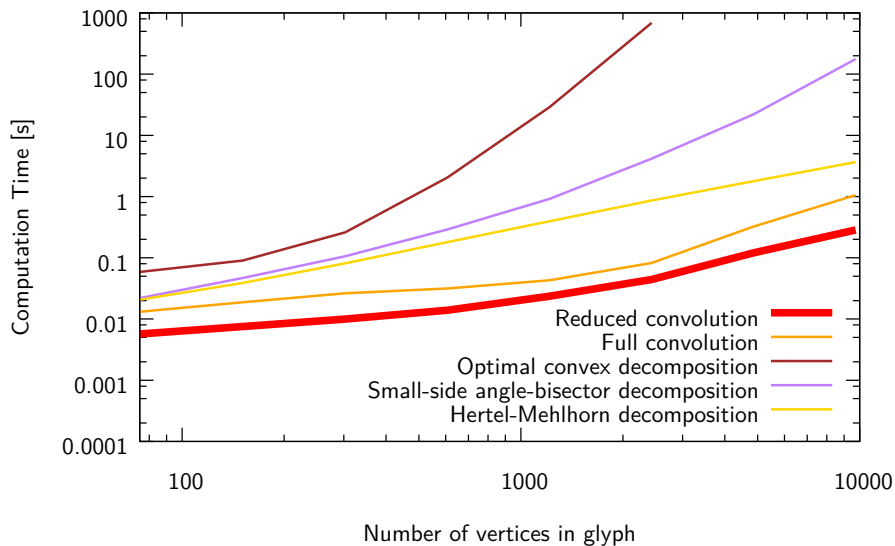


(75 vertices)

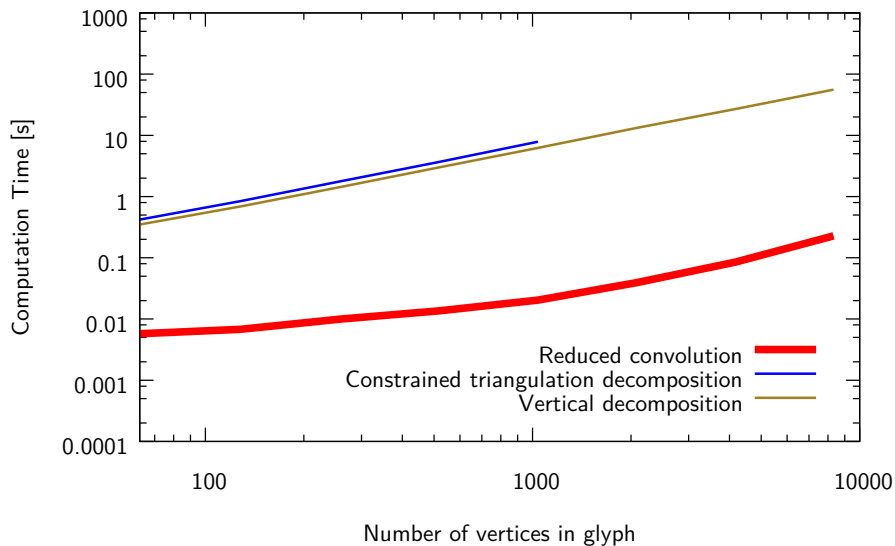


(8319 vertices)

# Benchmark: Glyph Offset (letter M)



# Benchmark: Glyph Offset (letter A)



## Contributions

- 1 Minkowski Sum of polygons with holes in CGAL
  - ▶ reduced convolution
  - ▶ two decomposition methods

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- ① Minkowski Sum of polygons with holes in CGAL
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- ② General input-level **hole filter** to reduce complexity

## Contributions

- 1 Minkowski Sum of polygons with holes in CGAL
  - ▶ reduced convolution
  - ▶ two decomposition methods
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Thanks!

